

Elimination Theory in Positive Characteristic

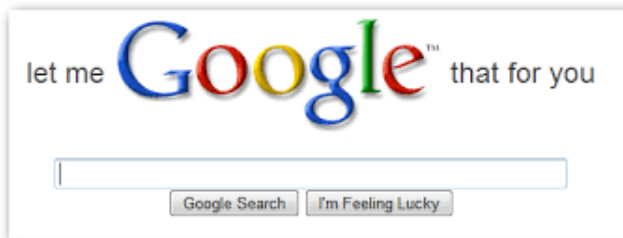
Carlos D'Andrea

June 23rd, 2016



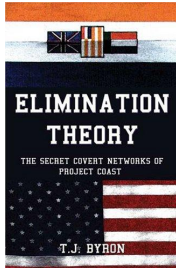
What is Elimination Theory?

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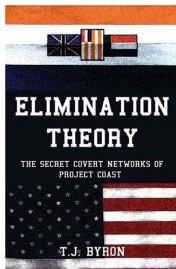


Elimination Theory

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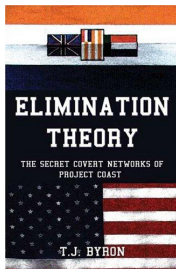
Elimination Theory



$\text{Res}(f(x), g(x)) =$

$$\begin{vmatrix} a_n & a_{n-1} & \cdots & a_1 & a_0 & 0 & \cdots & 0 \\ 0 & a_n & a_{n-1} & \cdots & a_1 & a_0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & \cdots & 0 & a_n & a_{n-1} & \cdots & a_1 & a_0 \\ b_n & b_{n-1} & \cdots & b_1 & b_0 & 0 & \cdots & 0 \\ 0 & b_n & b_{n-1} & \cdots & b_1 & b_0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & \cdots & 0 & b_n & b_{n-1} & \cdots & b_1 & b_0 \end{vmatrix}.$$

Elimination Theory



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Elimination Theory

Important for both **algorithmic** and **complexity** aspects of polynomial system solving



The Example: Determinants

Find “the condition” on
 $a_{10}, a_{11}, a_{20}, a_{21}$ so that the system

$$\begin{cases} a_{10}x_0 + a_{11}x_1 = 0 \\ a_{20}x_0 + a_{21}x_1 = 0 \end{cases}$$

has a solution different from $(0, 0)$

Elimination Theory

$$a_{10}x_0 + a_{11}x_1, a_{20}x_0 + a_{21}x_1$$

$$\in \mathbb{K}[a_{10}, a_{11}, a_{20}, a_{21}, x_0, x_1]$$

Elimination Theory

$$a_{10}x_0 + a_{11}x_1, a_{20}x_0 + a_{21}x_1$$

$$\in \mathbb{K}[a_{10}, a_{11}, a_{20}, a_{21}, x_0, x_1]$$

↓

$$a_{10}a_{21} - a_{20}a_{11} \in \mathbb{K}[a_{10}, a_{11}, a_{20}, a_{21}]$$

More general

Find “the condition” for the system

$$\left\{ \begin{array}{lcl} a_{10}x_0 + a_{11}x_1 + \dots + a_{1n}x_n & = & 0 \\ a_{20}x_0 + a_{21}x_1 + \dots + a_{2n}x_n & = & 0 \\ \vdots & & \vdots \\ a_{(n+1)0}x_0 + a_{(n+1)1}x_1 + \dots + a_{(n+1)n}x_n & = & 0 \end{array} \right.$$

More general

Find “the condition” for the system

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to have a solution different from
 $(0, 0, \dots, 0)$

Another more general

Let $d_1, d_2 \in \mathbb{N}$. Find “the condition”
for the system of polynomials

$$\begin{cases} a_{10}x_0^{d_1} + a_{11}x_0^{d_1-1}x_1 + \dots = 0 \\ a_{20}x_0^{d_2} + a_{21}x_0^{d_2-1}x_1 + \dots = 0 \end{cases}$$

to have a solution different from
 $(0, 0)$

More more more general...

Let $n \in \mathbb{N}$, and $d_1, \dots, d_{n+1} \in \mathbb{N}$, find the condition
for

$$\left\{ \begin{array}{lcl} \sum_{\alpha_0 + \dots + \alpha_n = d_1} a_{1, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} & = & 0 \\ \sum_{\alpha_0 + \dots + \alpha_n = d_2} a_{2, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} & = & 0 \\ & \vdots & \\ \sum_{\alpha_0 + \dots + \alpha_n = d_{n+1}} a_{n+1, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} & = & 0 \end{array} \right.$$

to have a solution different from $(0, 0, \dots, 0)$

Elimination: The general problem

For $\mathbf{a} = (a_1, \dots, a_N)$, $k, n \in \mathbb{N}$ let
 $f_1(\mathbf{a}, x_1, \dots, x_n), \dots, f_k(\mathbf{a}, x_1, \dots, x_n) \in$
 $\mathbb{K}[\mathbf{a}, x_1, \dots, x_n]$. Find conditions on $\mathbf{a}$ such that

$$\left\{ \begin{array}{lll} f_1(\mathbf{a}, x_1, \dots, x_n) & = & 0 \\ f_2(\mathbf{a}, x_1, \dots, x_n) & = & 0 \\ \vdots & \vdots & \vdots \\ f_k(\mathbf{a}, x_1, \dots, x_n) & = & 0 \end{array} \right.$$

has a solution

Solution?

- Depends on the ground field

Solution?

- Depends on the ground field
- There is not necessarily a “closed” condition

Solution?

- Depends on the ground field
- There is not necessarily a “closed” condition
- The computation of the conditions may be out of control

Easy example

$$\left\{ \begin{array}{lcl} a_{11}x_1 + \dots + a_{1n}x_n & = & 0 \\ a_{21}x_1 + \dots + a_{2n}x_n & = & 0 \\ \vdots & & \vdots \\ a_{kn}x_1 + \dots + a_{kn}x_n & = & 0 \end{array} \right.$$

with $k \geq n$

Easy example

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with $k \geq n$

Conditions: all maximal minors of $(a_{ij})_{1 \leq i \leq k, 1 \leq j \leq n}$
equal to zero

Another “easy” example

$$k = n = 1,$$

$$a_0 + a_1x_1 + a_2x_1^2 + \dots + a_dx_1^d = 0$$

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Conditions?

Geometry

$$V = \{(\mathbf{a}, x_1, \dots, x_n) : f_1(\mathbf{a}, x_1, \dots, x_n) = 0, \dots, f_k(\mathbf{a}, x_1, \dots, x_n) = 0\}$$

Geometry

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$$V \subset \mathbb{K}^N \times \mathbb{K}^n$$

$$\downarrow \pi_1$$

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$$\pi_1(V) \subset \mathbb{K}^N$$

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The set of conditions is $\pi_1(\mathbf{V})$, not necessarily described by zeroes of polynomials

Elimination Theorem

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$$\pi_1(V) = \{p_1(\mathbf{a}) = 0, \dots, p_\ell(\mathbf{a}) = 0\}$$

“The” Condition

$$V = \{(\mathbf{a}, x_0, x_1, \dots, x_n) : f_1(\mathbf{a}, x_0, x_1, \dots, x_n) = 0, \dots, f_{n+1}(\mathbf{a}, x_0, x_1, \dots, x_n) = 0\}$$

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Example 1

$$\left\{ \begin{array}{l} a_{00}x_0 + a_{01}x_1 + \dots + a_{0n}x_n = 0 \\ a_{10}x_0 + a_{11}x_1 + \dots + a_{1n}x_n = 0 \\ \qquad \qquad \qquad \vdots \qquad \qquad \qquad \vdots \qquad \vdots \\ a_{n0}x_0 + a_{n1}x_1 + \dots + a_{nn}x_n = 0 \end{array} \right.$$

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$$p_1(a) = \det(a_{ij})$$

Example 2

$$\begin{cases} f_1 = a_{10}x_0^{d_1} + a_{11}x_0^{d_1-1}x_1 + \dots + a_{1d_1}x_1^{d_1} \\ f_2 = a_{20}x_0^{d_2} + a_{21}x_0^{d_2-1}x_1 + \dots + a_{2d_2}x_1^{d_2} \end{cases}$$

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$$\text{Res}(f_1, f_2) = \det \begin{pmatrix} a_{10} & a_{11} & \dots & a_{1d_1} & 0 & \dots & 0 \\ 0 & a_{10} & \dots & a_{1d_1-1} & a_{1d_1} & \dots & 0 \\ \vdots & \vdots & \ddots & \dots & \dots & \ddots & \vdots \\ 0 & 0 & \dots & a_{10} & \dots & \dots & a_{1d_1} \\ a_{20} & a_{21} & \dots & a_{2d_2} & 0 & \dots & 0 \\ 0 & a_{20} & \dots & a_{2d_2-1} & a_{2d_2} & \dots & 0 \\ \vdots & \vdots & \ddots & \dots & \dots & \ddots & \vdots \\ 0 & 0 & \dots & a_{20} & \dots & \dots & a_{2d_2} \end{pmatrix}$$

Example 3

$$\left\{ \begin{array}{l} f_1 = \sum_{\alpha_0 + \dots + \alpha_n = d_1} a_{1, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} \\ f_2 = \sum_{\alpha_0 + \dots + \alpha_n = d_2} a_{2, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} \\ \vdots \\ f_{n+1} = \sum_{\alpha_0 + \dots + \alpha_n = d_{n+1}} a_{n+1, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} \end{array} \right.$$

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$$\boxed{\text{Res}(f_1, f_2, \dots, f_{n+1})}$$

The “elimination” dictionary

Linear

Polynomial

The “elimination” dictionary

Linear

Gauss elimination

Polynomial

Gröbner Bases

The “elimination” dictionary

Linear

Gauss elimination
Triangulation

Polynomial

Gröbner Bases
Triangular systems

The “elimination” dictionary

Linear

Gauss elimination

Triangulation

Determinants

Polynomial

Gröbner Bases

Triangular systems

Resultants

The “elimination” dictionary

Linear

Gauss elimination

Triangulation

Determinants

Cramer’s rule

...

Polynomial

Gröbner Bases

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Resultants

u-resultants

...

Positive characteristic

- The elimination theorems work independently of the characteristic of $\mathbb{K}$

Positive characteristic

- The elimination theorems work independently of the characteristic of $\mathbb{K}$
- We want to consider this scenario:
 $\mathbb{Z} \rightarrow \mathbb{F}_p$ for different primes $p \in \mathbb{Z}$

Warming-up example

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- 2 different roots if $p = 5$
- 3 different roots if $p \notin \{2, 3, 5\}$

Warming-up example

$f = 15x^3 + x^2 - 5x + 1 \in \mathbb{Z}[x]$ has

- 1 triple root if $p = 2$
- 1 double root if $p = 3$
- 2 different roots if $p = 5$
- 3 different roots if $p \notin \{2, 3, 5\}$
- 3 different roots over $\overline{\mathbb{Q}}$

Why do we care?

- Local-Global principles

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- Chinese Remainder like methods

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Local-Global Elimination Challenge

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Given $f_1, \dots, f_k \in \mathbb{Z}[x_1, \dots, x_n]$
describe $V_p(f_1, \dots, f_k) \subset \overline{\mathbb{F}_p}^n$

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(dimension, degree, height,
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■ for a given prime $p \in \mathbb{Z}$

Local-Global Elimination Challenge

Given $f_1, \dots, f_k \in \mathbb{Z}[x_1, \dots, x_n]$
describe $V_p(f_1, \dots, f_k) \subset \overline{\mathbb{F}}_p^n$
(dimension, degree, height,
irreducible components...)

- for a given prime $p \in \mathbb{Z}$
- for **all** prime $p \in \mathbb{Z}$
- Compare with $V_{\overline{\mathbb{Q}}}(f_1, \dots, f_k) \subset \overline{\mathbb{Q}}^n$

Interesting Analogy

Polynomial world
Systems over $\mathbb{Z}$

Interesting Analogy

Polynomial world

Systems over $\mathbb{Z}$

Linear world

Systems with parameters

$$\begin{bmatrix} -a & 1 & a \\ 1 & 0 & 0 \\ 0 & 1 & a \end{bmatrix}$$

Interesting-Nice fact

- There is a “generic” case

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- The generic case “is” the case over $\overline{\mathbb{Q}}$

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- There is a “generic” case
- The generic case “is” the case over $\overline{\mathbb{Q}}$
- We have size bounds for p to be generic

Example 1: Arithmetic Nullstellensatz

$$V_{\overline{\mathbb{Q}}}(f_1, \dots, f_k) = \emptyset$$

$$\iff$$

$$1 \in \langle f_1, \dots, f_k \rangle_{\mathbb{Q}}$$

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$$\iff \exists a \in \mathbb{Z} : a \in \langle f_1, \dots, f_k \rangle_{\mathbb{Z}}$$

and we have bounds for a

(D-Krick-Sombra 2013)

Corollary

If $p \nmid a$ then $V_p(f_1, \dots, f_k) = \emptyset$

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Nullstellensatz $\leftrightarrow$ “empty”
triangulation

$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 + \dots + a_{1n}x_n = b_1 \quad \dots (1)$$

$$a'_{22}x_2 + a'_{23}x_3 + \dots + a'_{2n}x_n = b'_2 \quad \dots (2)$$

$$a''_{33}x_3 + \dots + a''_{3n}x_n = b''_3 \quad \dots (3)$$

$$\vdots \quad \vdots \quad \vdots$$

$$a^{(n-1)}_{nn}x_n = b^{(n-1)}_n \quad \dots (n)$$

Example 2: zero dimensional systems

If $V_{\overline{\mathbb{Q}}}(f_1, \dots, f_k)$ is a set of T points,
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points

Example 2: zero dimensional systems

If $V_{\overline{\mathbb{Q}}}(f_1, \dots, f_k)$ is a set of T points,
then

$V_p(f_1, \dots, f_k)$ is also a set of T
points for $p \nmid a$ bounded

(D-Ostafe-Shparlinski-Sombra 2015)

What about the bad p 's?

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Anything can happen:

- Dimension change

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Anything can happen:

- Dimension change
- Degree change

What about the bad p 's?

$$\begin{bmatrix} -a & 1 & a \\ 1 & 0 & 0 \\ 0 & 1 & a \end{bmatrix}$$

Anything can happen:

- Dimension change
- Degree change
- Ramifications

What about the bad p 's?

$$\begin{bmatrix} -a & 1 & a \\ 1 & 0 & 0 \\ 0 & 1 & a \end{bmatrix}$$

Anything can happen:

- Dimension change
- Degree change
- Ramifications
- Embedded multiplicities

One interesting result

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Resultants modulo p

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Resultants modulo p

$$\text{Res}(f_1, f_2) = \det \begin{pmatrix} a_{10} & a_{11} & \dots & a_{1d_1} & 0 & \dots & 0 \\ 0 & a_{10} & \dots & a_{1d_1-1} & a_{1d_1} & \dots & 0 \\ \vdots & \vdots & \ddots & \dots & \dots & \ddots & \vdots \\ 0 & 0 & \dots & a_{10} & \dots & \dots & a_{1d_1} \\ a_{20} & a_{21} & \dots & a_{2d_2} & 0 & \dots & 0 \\ 0 & a_{20} & \dots & a_{2d_2-1} & a_{2d_2} & \dots & 0 \\ \vdots & \vdots & \ddots & \dots & \dots & \ddots & \vdots \\ 0 & 0 & \dots & a_{20} & \dots & \dots & a_{2d_2} \end{pmatrix}$$

Resultants modulo p

$$\text{Res}(f_1, f_2) = 0 \bmod p \iff \deg(\gcd(f_1 \bmod p, f_2 \bmod p)) > 0$$

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$$\text{Res}(f_1, f_2) = 0 \bmod p \iff \deg(\gcd(f_1 \bmod p, f_2 \bmod p)) > 0$$

$$p^{\deg(\gcd(f_1 \bmod p, f_2 \bmod p))} \mid \text{Res}(f_1, f_2)$$

(Gomez-Gutierrez-Ibeas-Sevilla 2009)

Questions

- What does this say about $V_p(f_1, f_2)$?

Questions

- What does this say about $V_p(f_1, f_2)$?
- Can you do it with more generality?

Vanishing of Resultantes modulo p

(D-Sombra 2016)

$$\left\{ \begin{array}{l} f_1 = \sum_{\alpha_0 + \dots + \alpha_n = d_1} a_{1, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} \\ f_2 = \sum_{\alpha_0 + \dots + \alpha_n = d_2} a_{2, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} \\ \vdots \\ f_{n+1} = \sum_{\alpha_0 + \dots + \alpha_n = d_{n+1}} a_{n+1, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} \end{array} \right.$$

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$$f_1, \dots, f_{n+1} \mathbb{Z}[\underline{x}] \implies \text{Res}(f_1, \dots, f_{n+1}) \in \mathbb{Z}$$

Known case

If $d_1 = d_2 = \dots = d_{n+1} = 1$, then

$$\text{Res}(f_1, \dots, f_{n+1}) = \det(a_{ij})_{1 \leq i, j \leq n+1}$$

A non trivial example

$$f_0 = a_{00}x_0 + a_{01}x_1 + a_{02}x_2$$

$$f_1 = a_{10}x_0 + a_{11}x_1 + a_{12}x_2$$

$$f_2 = a_{20}x_0^2 + a_{21}x_0x_1 + a_{22}x_0x_2 + a_{23}x_1^2 + a_{24}x_1x_2 + a_{25}x_2^2$$

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$$\begin{aligned} \text{Res}(f_0, f_1, f_2) = & a_{00}^2 a_{11}^2 a_{25} - a_{00}^2 a_{11} a_{12} a_{24} + a_{00}^2 a_{12}^2 a_{23} \\ & - 2a_{00} a_{01} a_{10} a_{11} a_{25} + a_{00} a_{01} a_{10} a_{12} a_{24} + a_{00} a_{01} a_{11} a_{12} a_{22} \\ & - a_{00} a_{01} a_{12}^2 a_{21} + a_{00} a_{02} a_{10} a_{11} a_{24} - 2a_{00} a_{02} a_{10} a_{12} a_{23} \\ & - a_{00} a_{02} a_{11}^2 a_{22} + a_{00} a_{02} a_{11} a_{12} a_{21} + a_{01}^2 a_{10}^2 a_{25} \\ & - a_{01}^2 a_{10} a_{12} a_{22} + a_{01}^2 a_{12}^2 a_{20} - a_{01} a_{02} a_{10}^2 a_{24} \\ & + a_{01} a_{02} a_{10} a_{11} a_{22} + a_{01} a_{02} a_{10} a_{12} a_{21} - 2a_{01} a_{02} a_{11} a_{12} a_{20} \\ & + a_{02}^2 a_{10}^2 a_{23} - a_{02}^2 a_{10} a_{11} a_{21} + a_{02}^2 a_{11}^2 a_{20} \end{aligned}$$

Properties of $\text{Res}(f_1, \dots, f_{n+1})$

- It is irreducible

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- It is irreducible
- It is homogeneous in each group of variables, of degree $\frac{d_1 \cdot d_2 \cdot \dots \cdot d_{n+1}}{d_i}$
- It is “weighted” homogeneous with weight $d_1 \cdot d_2 \cdot \dots \cdot d_{n+1}$

Geometric Properties

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- $\text{Res}(f_1, \dots, f_{n+1}) = 0 \iff$
 $\exists \xi \in \mathbb{P}^n$ such that
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■ Poisson Formula:

$$\begin{aligned} & \text{Res}(f_1, \dots, f_{n+1}) \\ &= \\ & \text{Res}(f_1^0, \dots, f_n^0)^{d_{n+1}} \prod_{\xi \in V(f_1^1, \dots, f_n^1)} f_{n+1}(\xi) \end{aligned}$$

Resolution of systems of polynomials

$$P(u_0, u_i) = \text{Res}(u_i x_0 - u_0 x_i, f_1, \dots, f_n)$$

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can be used to compute the
coordinates of the (finite) roots of
the system

$$f_1 = 0, \dots, f_n = 0$$

Computation

$$\mathcal{R}(f_0, f_1, f_2) = \det \begin{bmatrix} -b_1 & -b_3 & 0 & a_1 & a_3 & 0 & 0 & 0 & 0 & 0 \\ -b_0 & -b_2 & 0 & a_0 & a_2 & 0 & 0 & 0 & 0 & 0 \\ 0 & -b_1 & -b_3 & 0 & a_1 & a_3 & 0 & 0 & 0 & 0 \\ 0 & -b_0 & -b_2 & 0 & a_0 & a_2 & 0 & 0 & 0 & 0 \\ -c_4 & -c_5 & -c_8 & 0 & 0 & 0 & a_1 & 0 & a_3 & 0 \\ -c_1 & -c_3 & -c_7 & 0 & 0 & 0 & a_0 & a_1 & a_2 & a_3 \\ -c_0 & -c_2 & -c_6 & 0 & 0 & 0 & 0 & a_0 & 0 & a_2 \\ 0 & 0 & 0 & -c_4 & -c_5 & -c_8 & b_1 & 0 & b_3 & 0 \\ 0 & 0 & 0 & -c_1 & -c_3 & -c_7 & b_0 & b_1 & b_2 & b_3 \\ 0 & 0 & 0 & -c_0 & -c_2 & -c_6 & 0 & b_0 & 0 & b_2 \end{bmatrix}$$

Resultants modulo p

(D-Sombra 2016)

If $\dim \left(V_p(f_1, \dots, f_{n+1}) \right) \leq 0$

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$$\text{If } \dim \left(V_p(f_1, \dots, f_{n+1}) \right) \leq 0 \\ N_p := \deg \left(V_p(f_1, \dots, f_{n+1}) \right)$$

Resultants modulo p

(D-Sombra 2016)

If $\dim(V_p(f_1, \dots, f_{n+1})) \leq 0$
 $N_p := \deg(V_p(f_1, \dots, f_{n+1}))$

$$p^{N_p} \mid \text{Res}(f_1, \dots, f_{n+1})$$

Corollary

The factorization of the resultant
actually bounds the (finite) zeroes
modulo p !

$$p^{N_p} \mid \text{Res}(f_1, \dots, f_{n+1})$$

The Cantabrian Theorem revisited

$$\begin{aligned} &\deg(\gcd(f_1 \bmod p, f_2 \bmod p)) \\ &= \\ &N_p \\ &= \\ &\deg(V_p(f_1, f_2)) \end{aligned}$$

(Gomez-Gutierrez-Ibeas-Sevilla 2009)

Remarks

- Still the result works under the (generic) hypothesis of finiteness modulo p

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- the “gap” in the bound can be large
- Not a clear “algorithm” for deciding when $\dim(V_p(f_1, \dots, f_{n+1})) > 0$

Idea of our proof

- “Remove” all the zeroes from the infinite

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- Compute the dimension of the Nullspace of the determinantal matrix

Generalizations and Extensions

- One could get a refinement of the exponent by taking into account the valuation modulo p of the roots (Smirnov's Theorem)

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- One could get a refinement of the exponent by taking into account the valuation modulo p of the roots (Smirnov's Theorem)
- Slight generalization to *sparse resultants* under stronger hypothesis
- The result holds for any domain, for instance polynomials with coefficients in $R[y_1, \dots, y_l]$

Applications

- Finding points in varieties modulo p (Shparlinski)

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- Removing some “extraneous factors” in the Computation of the “Salmon Polynomial” (Busé-Chardin-D-Sombra-Weimann)

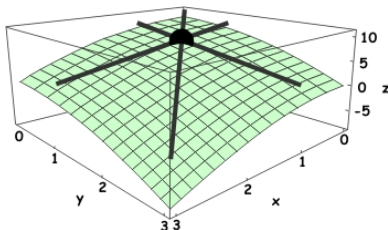
Computation of the Salmon Polynomial

(Busé-Chardin-D-Sombra-Weimann)

$$\mathbb{Z} \leftrightarrow \mathbb{C}[x, y, z]$$

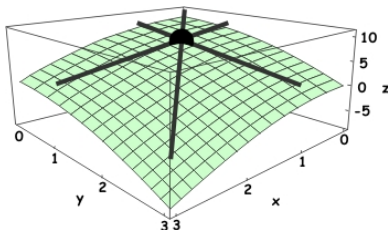
$$p \leftrightarrow f(x, y, z)$$

Salmon's polynomial



A point b in a surface $S \subset \mathbb{C}^3$ is called *flex* (or inflection) of S

Salmon's polynomial



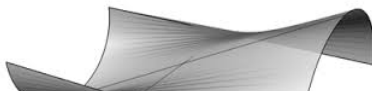
A point b in a surface $S \subset \mathbb{C}^3$ is called *flex* (or *inflection*) of S if there exists a line passing through b having contact order at least 3 with S

Theorem (Salmon, 1862)

If $S = V(f(x, y, z)) \subset \mathbb{C}^3$ of degree d , and not ruled, there is

$F_f(x, y, z) \in \mathbb{C}[x, y, z]$ of degree $\leq 11d - 24$ such that

$$\text{Flex}(S) = V(f(x, y, z), F_f(x, y, z))$$



Computing $F_f(x, y, z)$

$$f((x, y, z) + t(u, v, w)) =$$

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$$\begin{aligned} f((x, y, z) + t(u, v, w)) = \\ f(x, y, z) + t f_1(x, y, z; u, v, w) + \\ t^2 f_2(x, y, z; u, v, w) + \\ t^3 f_3(x, y, z; u, v, w) + \mathcal{O}(t^4) \end{aligned}$$

Computing $F_f(x, y, z)$

The “candidate” for $F_f(x, y, z)$ should be the resultant in (u, v, w) of

- $f_1(x, y, z; u, v, w)$
- $f_2(x, y, z; u, v, w)$
- $f_3(x, y, z; u, v, w)$

Janos Kollár (arxiv: 2014)

“I get a polynomial of degree $11d - 18$. Salmon claims that in fact the degree should be $11d - 24$. I have not checked this”

Terence Tao (blog, 2014)

“The original proof of the Cayley-Salmon theorem, dating back to at least 1915, is not easily accessible and not written in modern language”

Our Result

(Busé-Chardin-D-Sombra-Weimann)

Working modulo $f(x, y, z)$, the
resultant has a spurious factor of
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$$11d - 18 - 6 = 11d - 24$$

Thanks!

