

# Homological Processing of Biomedical digital images: automation and certification<sup>1</sup>

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- 1 Motivation
- 2 Automating the process
- 3 Main problems
- 4 Conclusions and further work

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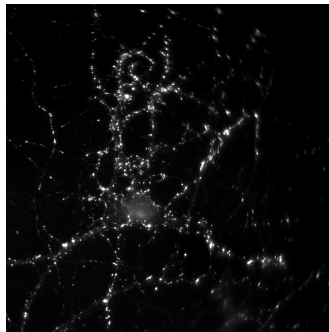
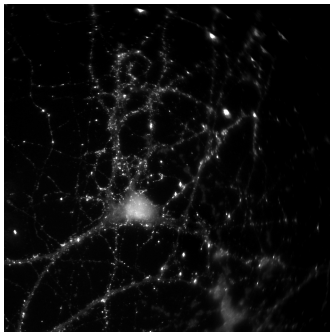
- 1 Motivation
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# Motivation: Synapses counting

- *Synapses* are the points of connection between neurons
- *Relevance*: Computational capabilities of the brain
- The different number of synapses may be an important asset in the treatment of neurological diseases

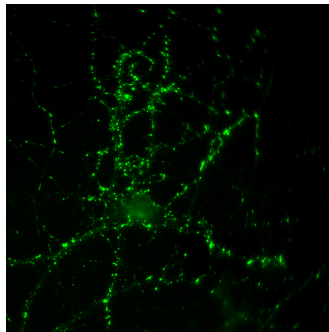
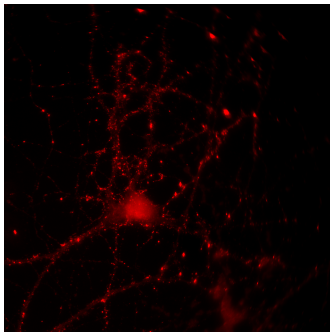
# Manual processing to count synapses

- Apply two different antibody markers, bassoon and synapsin



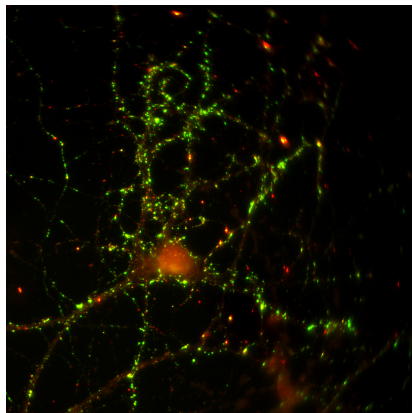
## Manual processing to count synapses

- Process the images in order to count the synapses (*ImageJ*)



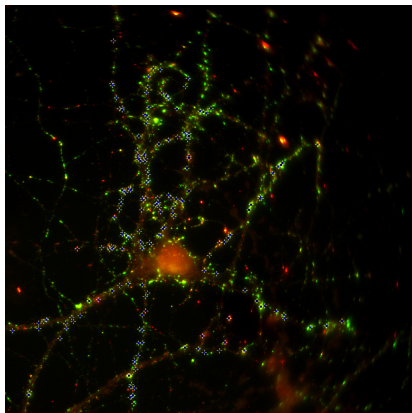
# Manual processing to count synapses

- Overlap both images



## Manual processing to count synapses

- The synapses are manually counted one by one





# Goal to count synapses

## Goal

Provide a reliable and automatic method for counting synapses in a neuron

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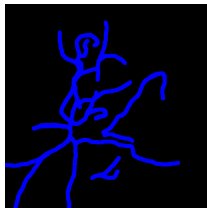
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# Count synapses automatically

- New ImageJ plugin called *SynapCountJ*
- Improve the interaction with the *ImageJ* system

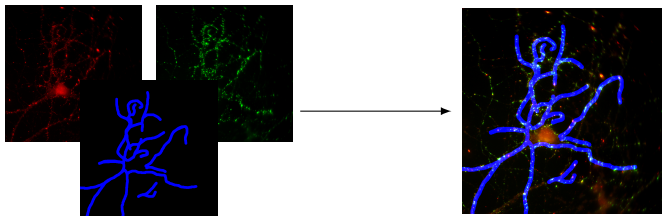
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- New ImageJ plugin called *SynapCountJ*
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- Steps
  - Determine the neuron morphology from one of those pictures (*NeuronJ plugin*)



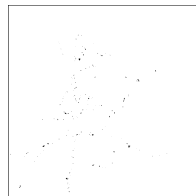
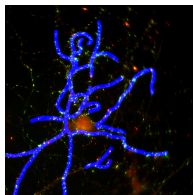
# Count synapses automatically

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  - Overlap the images with the two markers with the one with the structure (*SynapCountJ*)



# Count synapses automatically

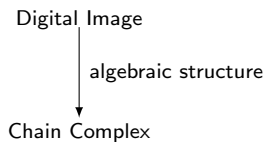
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- Improve the interaction with the *ImageJ* system
- Steps
  - Determine the neuron morphology from one of those pictures (*NeuronJ plugin*)
  - Overlap the images with the two markers with the one with the structure (*SynapCountJ*)
  - Invert the colors to show the synapses as black points



# The method

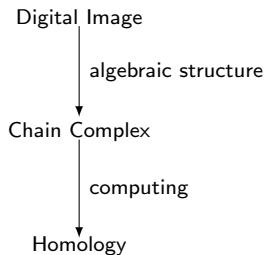
Digital Image

# The method

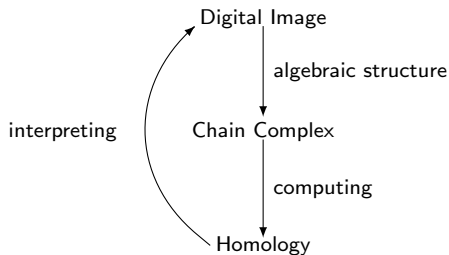




# The method



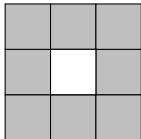
# The method



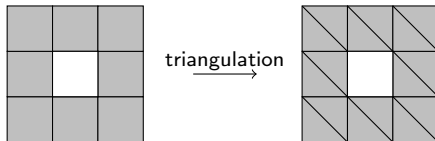
# Image to Chain Complex



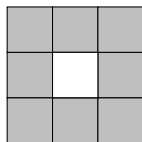
# Image to Chain Complex



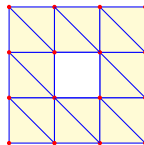
# Image to Chain Complex



# Image to Chain Complex



triangulation  $\longrightarrow$



$$C_0 = \mathbb{Z}[\text{vertices}]$$

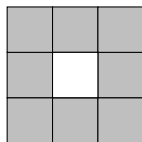
$\uparrow d_1$

$$C_1 = \mathbb{Z}[\text{edges}]$$

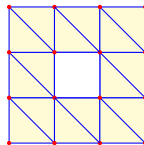
$\uparrow d_2$

$$C_2 = \mathbb{Z}[\text{triangles}]$$

# Image to Chain Complex



triangulation  $\longrightarrow$



$\longrightarrow$

$$C_0 = \mathbb{Z}[\text{vertices}]$$

$\uparrow d_1$

$$C_1 = \mathbb{Z}[\text{edges}]$$

$\uparrow d_2$

$$C_2 = \mathbb{Z}[\text{triangles}]$$

$$0 \leftarrow \mathbb{Z}^{16} \xleftarrow{d_1} \mathbb{Z}^{32} \xleftarrow{d_2} \mathbb{Z}^{16} \leftarrow 0$$

# Compute Homology



- Problem of diagonalizing matrices
- Compute the Smith Normal Form



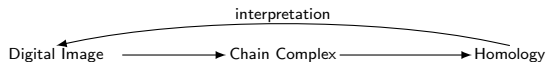
# Interpretation from Homology to Image



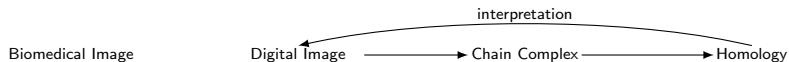
- $H_0$  measures the number of connected components
- $H_1$  measures the number of holes

In our case,  $H_0$  counts the number of synapses

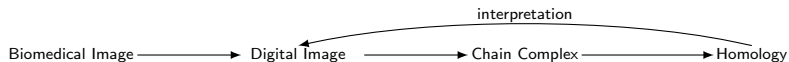
# General method



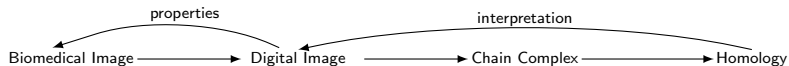
# General method



# General method



# General method



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- 2 Automating the process
  - Preprocessing the image
  - Algebraic Topology for digital images analysis
  - General method
- 3 Main problems
  - Reduce the size: Discrete Morse theory
  - Safety of the results: Certification of the programs
- 4 Conclusions and further work

# Problems

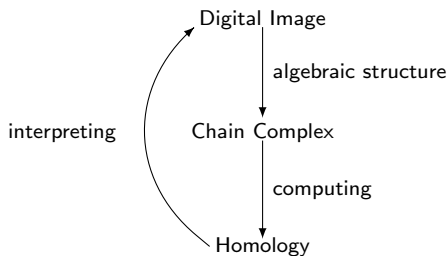
- Size of the images
- Safety of the results

# Problems

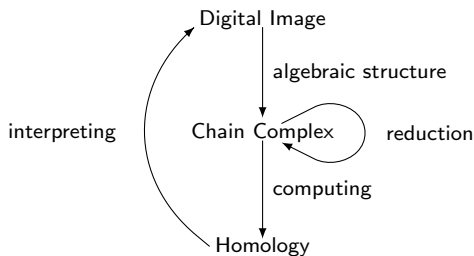
- Size of the images  $\rightarrow$  Discrete Morse theory
- Safety of the results  $\rightarrow$  Certification of the programs



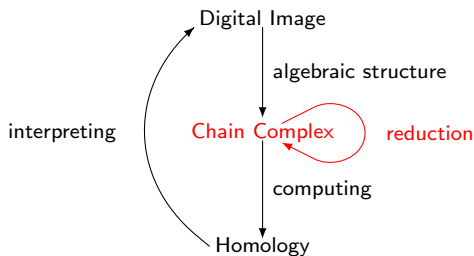
# The method



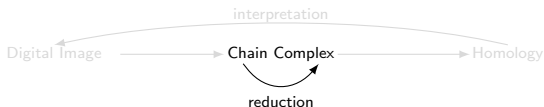
# The method



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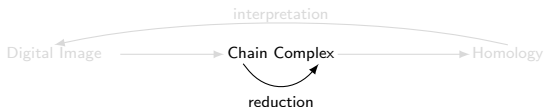


# Reduction of chain complex

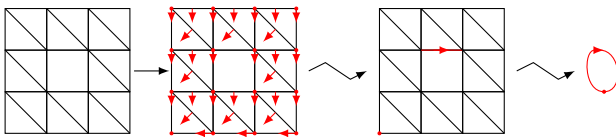


- Reduce information keeping the homological properties
- Discrete Morse Theory
  - Vector fields are a tool to cancel “useless” information

# Reduction of chain complex

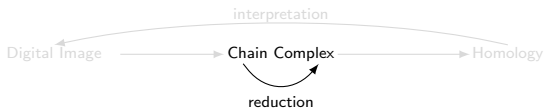


- Reduce information keeping the homological properties
- Discrete Morse Theory
  - Vector fields are a tool to cancel “useless” information

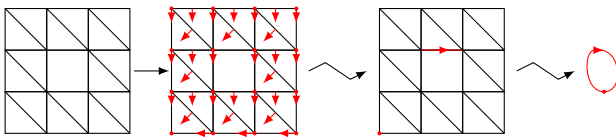


$$0 \leftarrow \mathbb{Z}^{16} \leftarrow \mathbb{Z}^{32} \leftarrow \mathbb{Z}^{16} \leftarrow 0$$

# Reduction of chain complex



- Reduce information keeping the homological properties
- Discrete Morse Theory
  - Vector fields are a tool to cancel “useless” information



$$0 \leftarrow \mathbb{Z} \leftarrow \mathbb{Z} \leftarrow 0 \leftarrow 0$$

# Discrete Morse Theory

## Definition

Let  $C_* = (C_p, d_p)_{p \in \mathbb{Z}}$  be a free chain complex with distinguished  $\mathbb{Z}$ -basis  $\beta_p \subset C_p$ . A  $(p-1)$ -cell  $\sigma$  is a *face* of a  $p$ -cell  $\tau$  if the coefficient of  $\sigma$  in  $d\tau$  is non-null. It is a *regular face* if this coefficient is  $+1$  or  $-1$

## Definition

A *discrete vector field* on  $C_*$  is a collection of pairs  $V = \{(\sigma_i, \tau_i)\}_{i \in \beta}$  satisfying the conditions:

- 1 Every  $\sigma_i$  is some element of  $\beta_p$ , in which case the other corresponding component  $\tau_i \in \beta_{p+1}$ . The degree  $p$  depends on  $i$  and in general is not constant
- 2 Every component  $\sigma_i$  is a *regular face* of the corresponding component  $\tau_i$
- 3 A generator of  $C_*$  appears at most one time in  $V$

# Discrete Morse Theory

## Definition

A  $V$ -path of degree  $p$  is a sequence  $\pi = ((\sigma_{i_k}, \tau_{i_k}))_{0 \leq k < m}$  satisfying:

- 1 Every pair  $((\sigma_{i_k}, \tau_{i_k}))$  is a component of  $V$  and the cell  $\tau_{i_k}$  is a  $p$ -cell
- 2 For every  $0 < k < m$ , the component  $\sigma_{i_k}$  is a face of  $\tau_{i_{k-1}}$ , non necessarily regular, but different from  $\sigma_{i_{k-1}}$

## Definition

A discrete vector field  $V$  is admissible if for every  $p \in \mathbb{Z}$ , a function  $\lambda_p : \beta_p \rightarrow \mathbb{Z}$  is provided satisfying the property: every  $V$ -path starting from  $\sigma \in \beta_p$  has a length bounded by  $\lambda_p(\sigma)$



# Discrete Morse Theory

## Definition

A cell  $\chi$  which does not appear in a discrete vector field  $V = \{(\sigma_i, \tau_i)\}_{i \in \beta}$  is called a *critical cell*

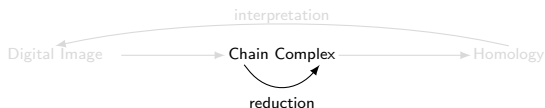
## Vector-Field Reduction Theorem

Let  $C_* = (C_p, d_p \beta_p)_p$  be a free chain complex and  $V = \{(\sigma_i, \beta_i)\}_{i \in \beta}$  be an admissible discrete vector field on  $C_*$ . Then the vector field  $V$  defines a canonical reduction  $\rho = (f, g, h) : (C_p, d_p) \implies (C_p^c, d_p')$  where  $C_p^c = \mathbb{Z} [\beta_p^c]$  is the free  $\mathbb{Z}$ -module generated by the critical  $p$ -cells



A. Romero and F. Sergeraert. Discrete Vector Fields and Fundamental Algebraic Topology, 2010. <http://arxiv.org/abs/1005.5685v1>.

# Discrete vector field over matrices

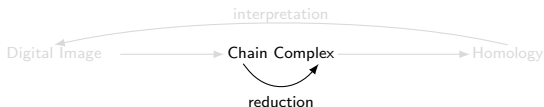


- 2D-images
- Chain complex associated with an image is finite

$$0 \leftarrow C_0 \xleftarrow{d_1} C_1 \xleftarrow{d_2} C_2 \leftarrow 0$$

- Differential maps can be represented by integer matrices
- Reduction chain complex  $\rightarrow$  Reduction matrices

# Discrete vector field over matrices

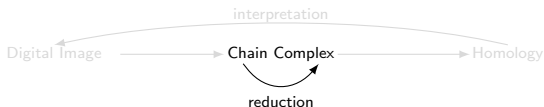


## Definition

A *vector field*  $V$  for a matrix  $M \in \text{Mat}_{m,n}(\mathbb{Z})$  is a set of integer pairs  $\{(a_i, b_i)\}_i$  satisfying these conditions:

- 1  $1 \leq a_i \leq m$  and  $1 \leq b_i \leq n$
- 2 The entry  $M[a_i, b_i]$  is  $\pm 1$
- 3 The indices  $a_i$  (respectively  $b_i$ ) are pairwise different

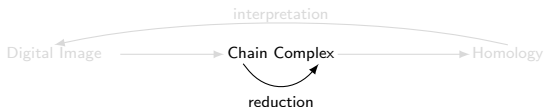
# Reduction of Chain Complex



## Goal

- Let  $M_n$  be a finite matrix which represents the differential map  $d_n$  of  $C_*$ 
  - Compute an admissible discrete vector field  $V$  from  $M_n$
  - Obtain a new matrix  $\hat{M}_n$  from  $M_n$  and  $V$

# Reduction of Chain Complex

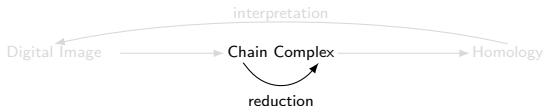


## Goal

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  - Compute an admissible discrete vector field  $V$  from  $M_n$
  - Obtain a new matrix  $\hat{M}_n$  from  $M_n$  and  $V$

In our case, we have to reduce two matrices  $M_1$  and  $M_2$ .  
Compute the homology groups of  $C_*$  with  $\hat{M}_1$  and  $\hat{M}_2$  can be much faster.

# Reduction of Chain Complex



- Implemented in Haskell

## Algorithm 1

*Input:* an integer matrix  $M_n$

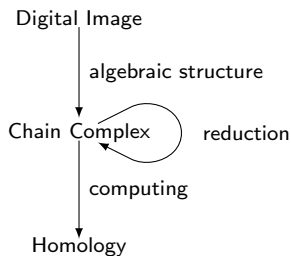
*Output:* an admissible discrete vector field  $V$

## Algorithm 2

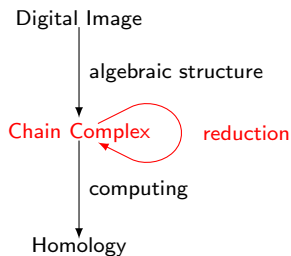
*Input:* an integer matrix  $M_n$

*Output:* a reduced matrix  $\hat{M}_n$

# The method



# The method





# Coq/SSReflect

- Coq
  - Theorem prover assistant
  - High logic order
- SSReflect
  - Extension of Coq
  - Introduce new tactics and libraries
  - Used to formalize of the Four Colour Theorem

# Reduction of chain complex



- Steps

- 1 Translate our *Haskell* code into the *Coq* language
- 2 Define the test functions to specify the properties which our programs must satisfy
- 3 State and prove the lemmas which ensure the correctness of our programs

# Reduction of chain complex



- Steps

- ① Translate our *Haskell* code into the *Coq* language
- ② Define the test functions to specify the properties which our programs must satisfy
- ③ State and prove the lemmas which ensure the correctness of our programs

**Example:** Let  $M$  be an integer matrix, `(vectorCvd M)` builds an admissible discrete vector field

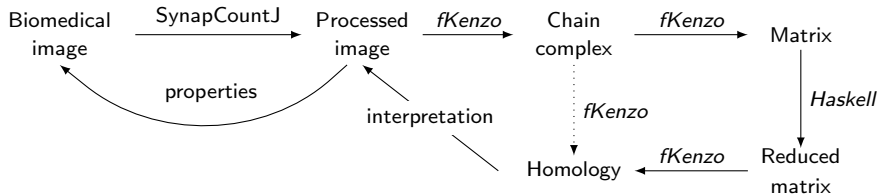
.....  
Lemma admissible-vf:

forall M, (int-matrix M) -> (admissible (vectorCvd M))  
.....

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# Conclusions



- Methodology to study Biomedical images
- Programs partially verified with Theorem Prover tools
- Application to count synapses

## Further work

- Verification of our *Haskell* programs by means of *Coq/SSReflect* is still an ongoing work
- Verification of Smith Normal Form of a matrix
- Find other applications of our homological tools in the Biomedical imaging context

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